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# Exact Pareto-Optimal Coordination of Two Translating Polygonal Robots on an Acyclic Roadmap

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# Introduction

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We want to study coordination strategies for robots in a shared workspace. We allow individual robots to have separate performance measures.

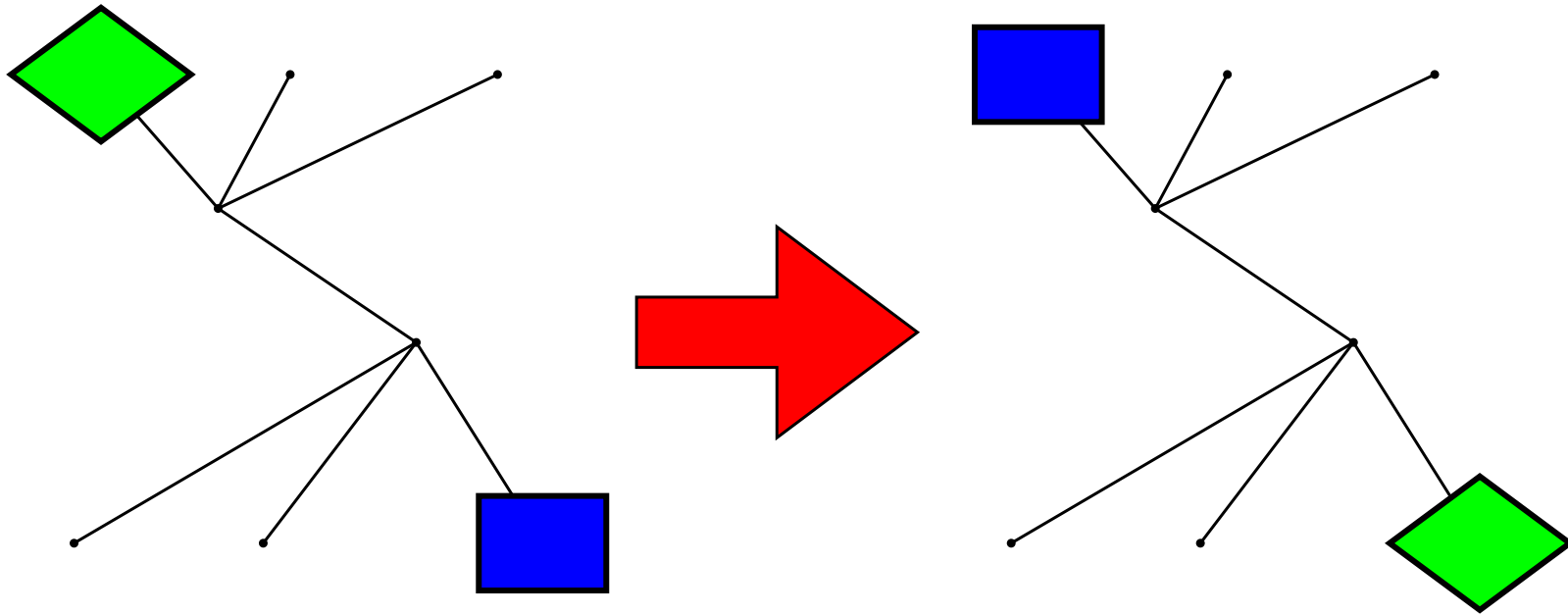
*Problem:* Find collision-free motion strategies that are optimal in a multi-objective sense.

*Applications:*

- AGVs in a factory setting.
- General multiple robot coordination.

# Introduction

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# Overview

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- Related Work
- Problem Statement
- Robots on Fixed Paths
- Generalization to Roadmaps
- Examples

# Related Work

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## *Centralized Methods*

Schwartz and Sharir, 1983

Ardema and Skowronski, 1991

Barraquand and Latombe, 1991

## *Decentralized (“Fixed Path”) Methods*

Erdmann and Lozano-Perez, 1986

Akella and Hutchinson, 2002

Siméon, Leroy and Laumond 2002

Peng and Akella, 2003

## *Roadmap Methods*

LaValle and Hutchinson, 1998

# Problem Statement

Two polygonal robots  $\mathcal{R}_1$  and  $\mathcal{R}_2$  translating in the plane.

- Robots move on roadmaps  $\mathcal{G}_1$  and  $\mathcal{G}_2$  of piecewise-linear paths.
- Initial and goal configurations  $X_i^{init}, X_i^{goal} \in \mathcal{G}_i$ .
- Allow instantaneous changes in speed.

*Objective:* Find a continuous collision-free path

$$C : [0, 1] \rightarrow \mathcal{G}_1 \times \mathcal{G}_2$$

from  $(X_1^{init}, X_2^{init})$  to  $(X_1^{goal}, X_2^{goal})$ .

# Optimality

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For a given coordination, each robot has a cost function:

$$J = (J_1, J_2)$$

One approach is to choose a scalarization function  $f : \mathbb{R}^2 \rightarrow \mathbb{R}$  and optimize  $f(J)$ .

- Scalarizing may omit interesting solutions.
- Priorities may change across multiple queries.
- Better to find a small set of good candidate solutions.

# Pareto Optimality

Rather than choosing any particular scalarization, we find the set of **Pareto optimal** solutions.

- Create equivalence classes of paths with identical costs.

$$C \sim C' := J(C) = J(C')$$

- Define a partial order on equivalence classes:

$$[C] \leq [C'] := J_1(C) \leq J_1(C') \wedge J_2(C) \leq J_2(C')$$

- Pareto optima are the minima in this partial order.



# Coordination Space

O'Donnell and Lozano-Perez, 1989

We want to find a path through the *coordination space*  $\mathcal{G}_1 \times \mathcal{G}_2$ .

- Obstacle regions where  $\mathcal{R}_1$  collides with  $\mathcal{R}_2$ .

The slope of this curve determines the velocity of each robot.

- Slope  $\geq 1$ :  $\mathcal{R}_2$  at full speed
- Slope  $\leq 1$ :  $\mathcal{R}_1$  at full speed.

Time to execute a segment is its  $L^\infty$  length.

# Optima for Fixed Paths

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*Lemma:* Every Pareto-optimal coordination class contains a coordination composed of segments of the visibility graph of the obstacle set, plus possibly a “full speed completion.”

*Proof Ideas:*

- Given any path, “shorten” it until it’s constrained by obstacle vertices.
- After moving past the last obstruction, both robots should move at full speed to their goals.

# Optima for Fixed Paths

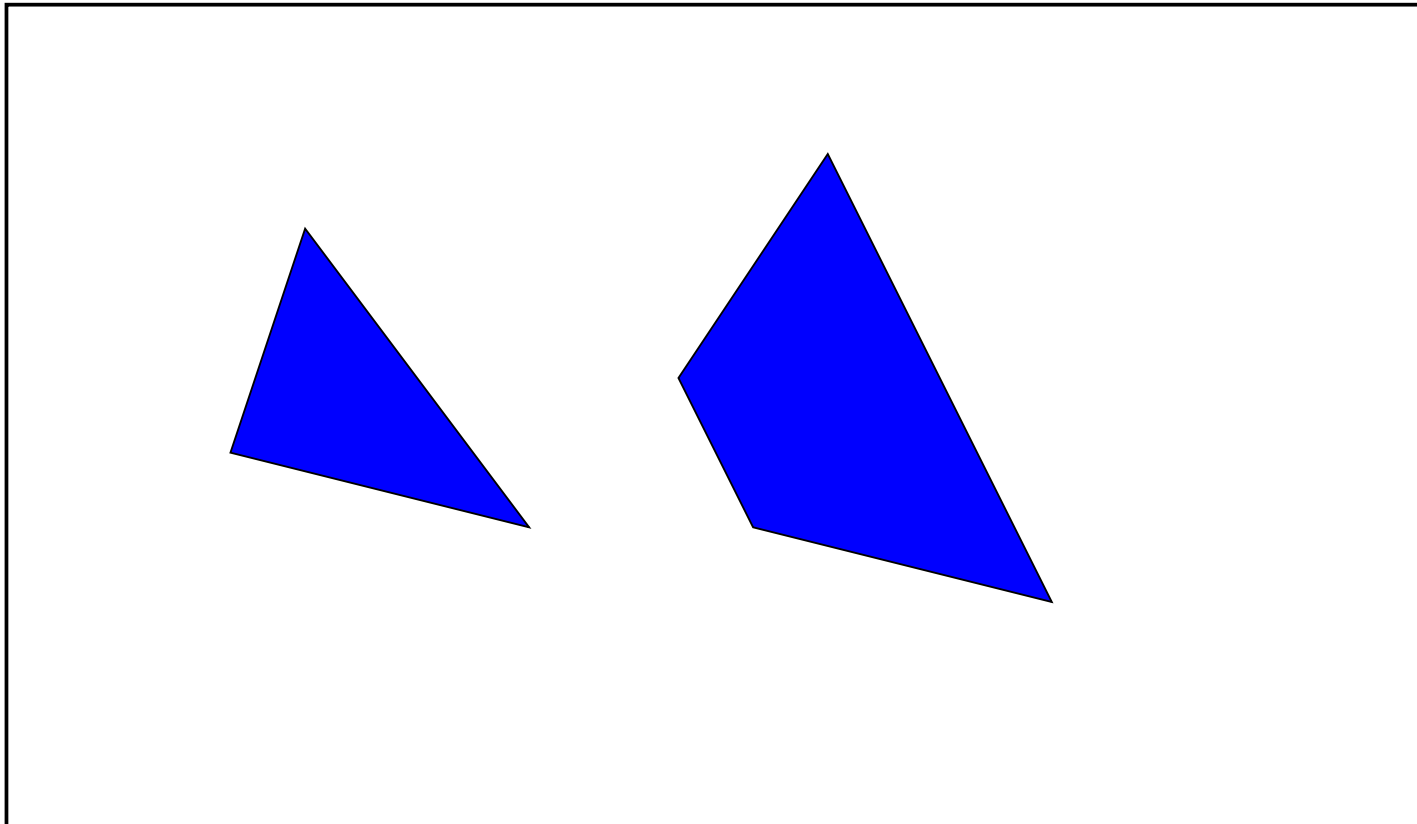
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## *Algorithm:*

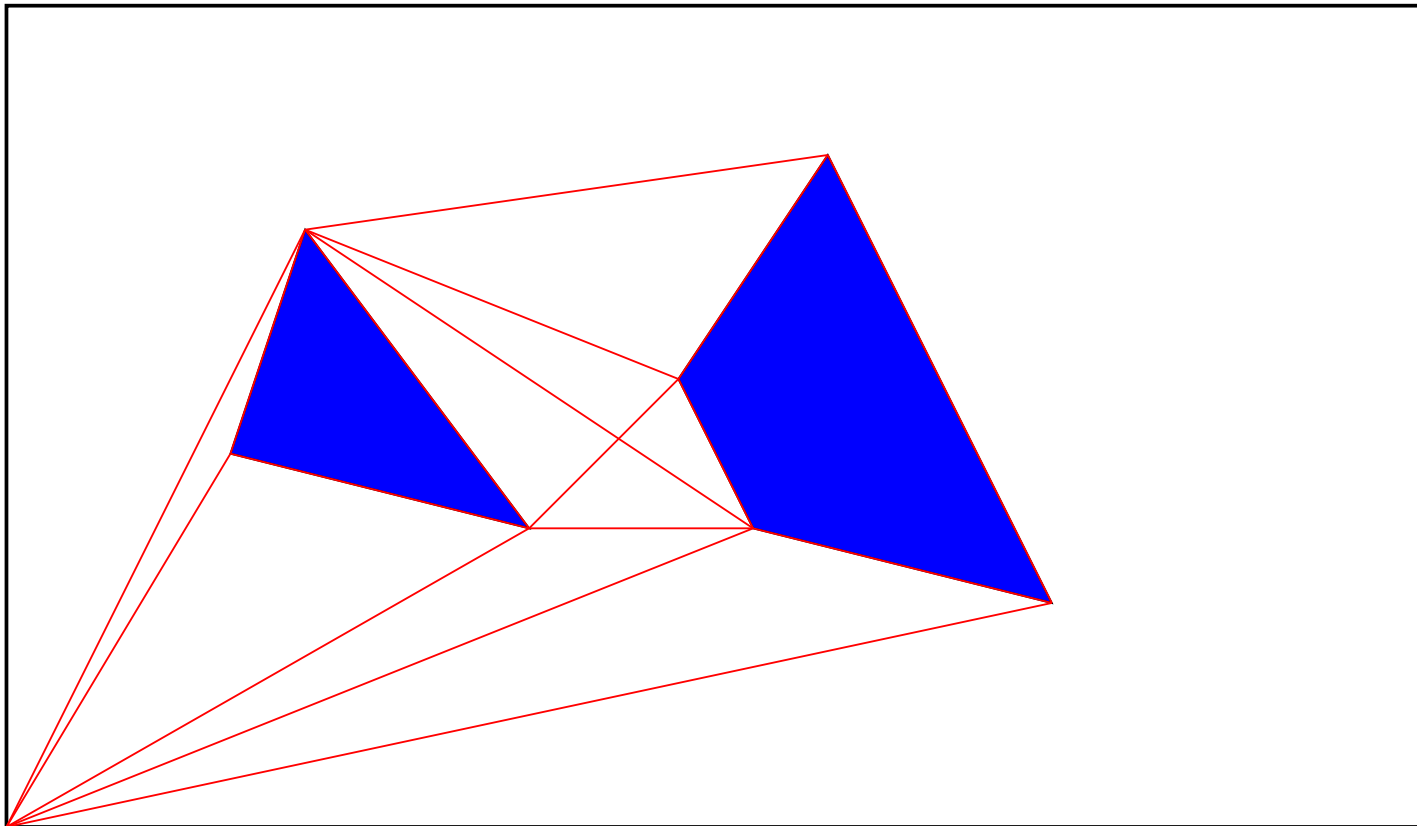
- Compute the obstacle set.
- Find the visibility graph of obstacle set.
- Add a full-speed completion from each vertex for which this is possible.
- Use Dijkstra's algorithm to extract a set of candidate solutions.
  - Candidate = Shortest path to obstacle vertex + full-speed completion.
- Use direct comparisons to eliminate candidates that are not Pareto optima.

# Fixed Path Example

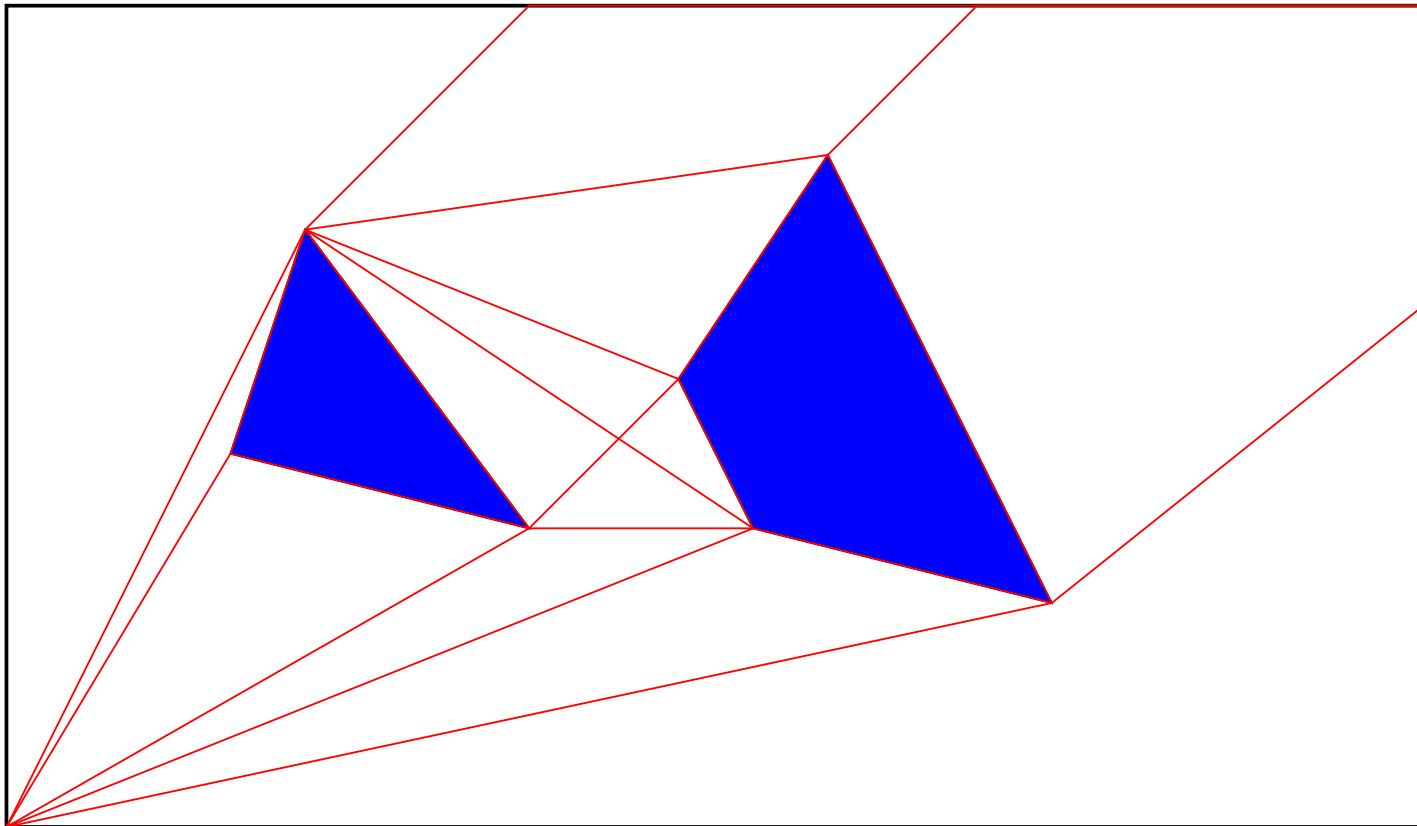
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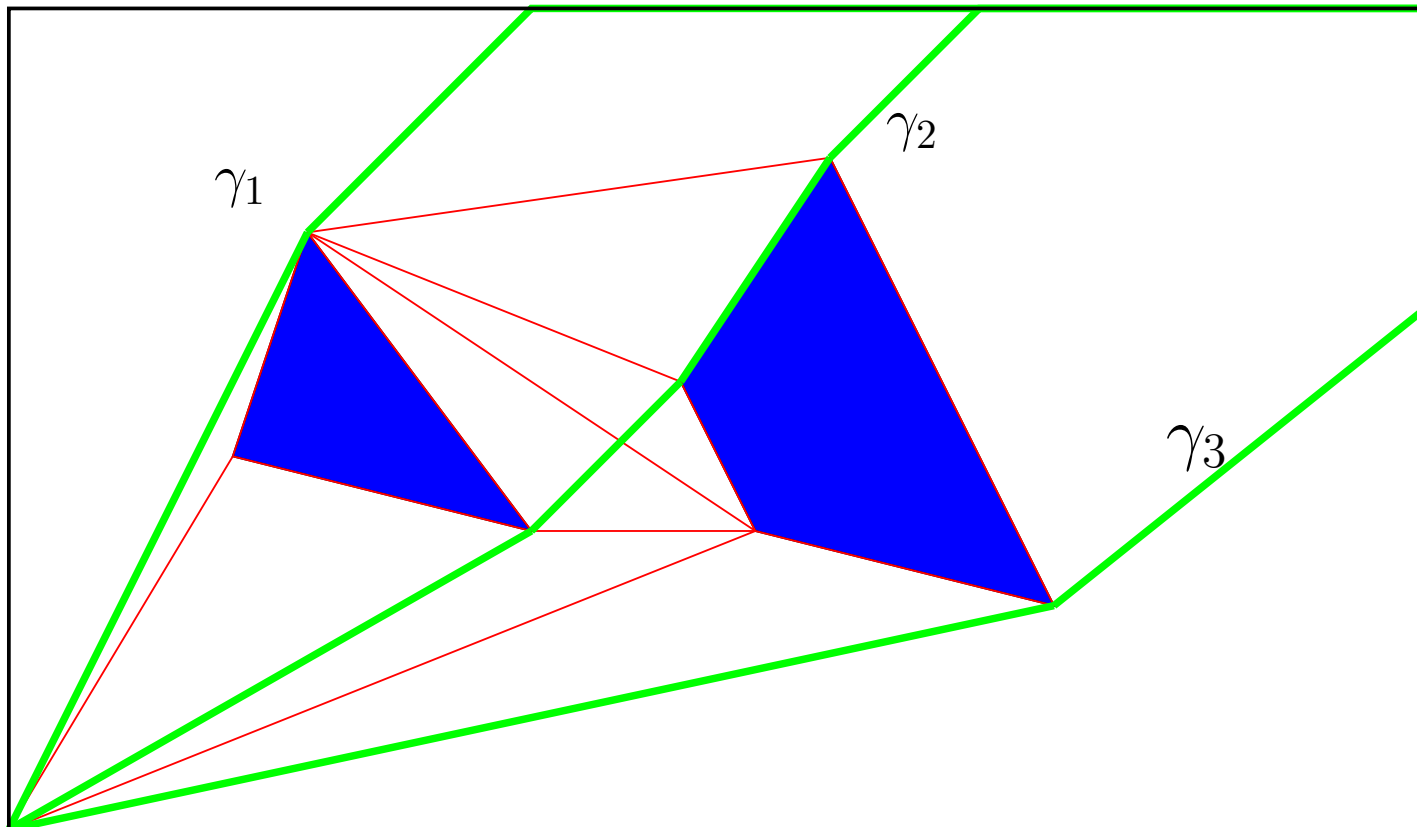
# Fixed Path Example



# Fixed Path Example



# Fixed Path Example

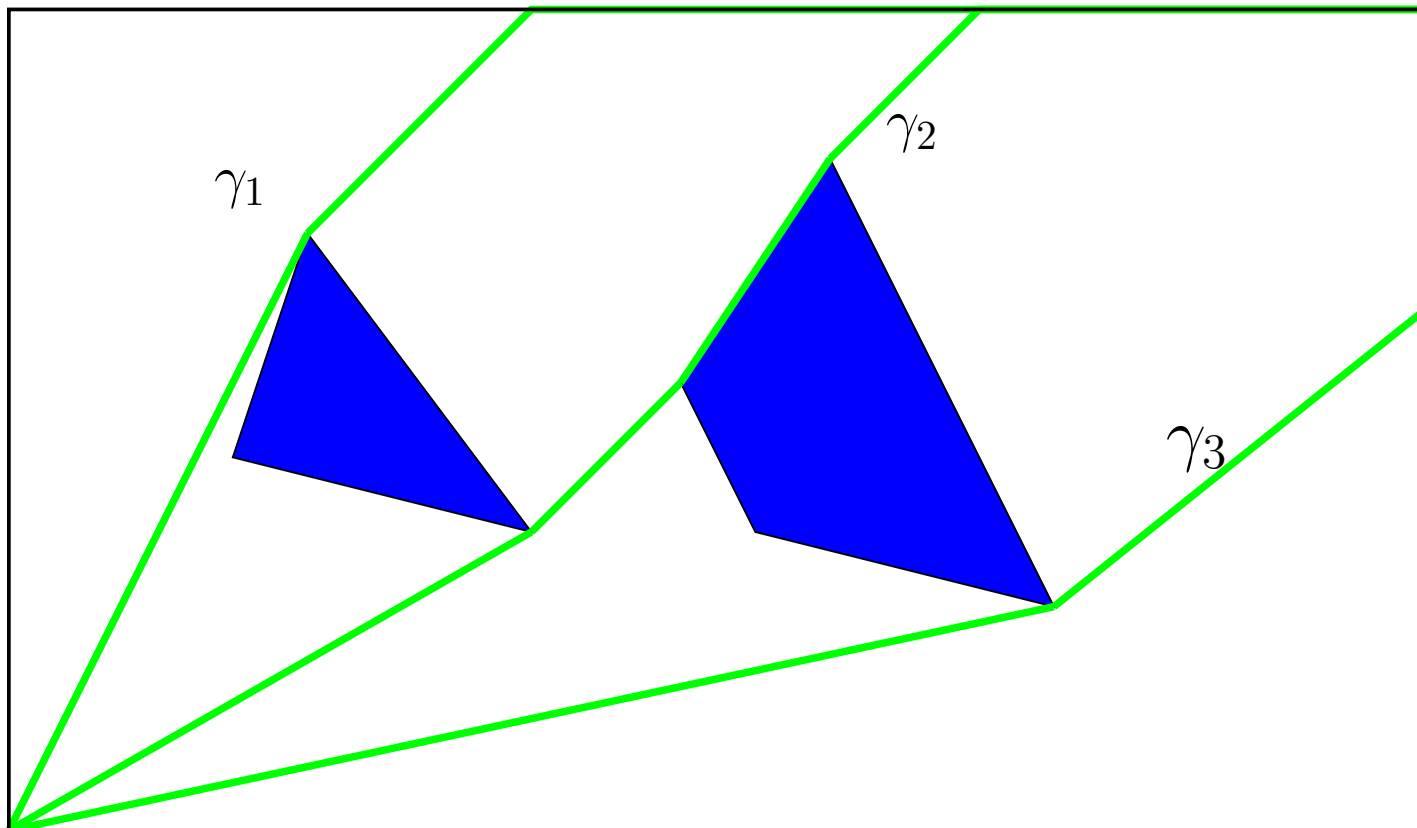


$$J(\gamma_1) = (23, 11)$$

$$J(\gamma_2) = (21, 15)$$

$$J(\gamma_3) = (19, 25)$$

# Fixed Path Example



$$J(\gamma_1) = (23, 11)$$

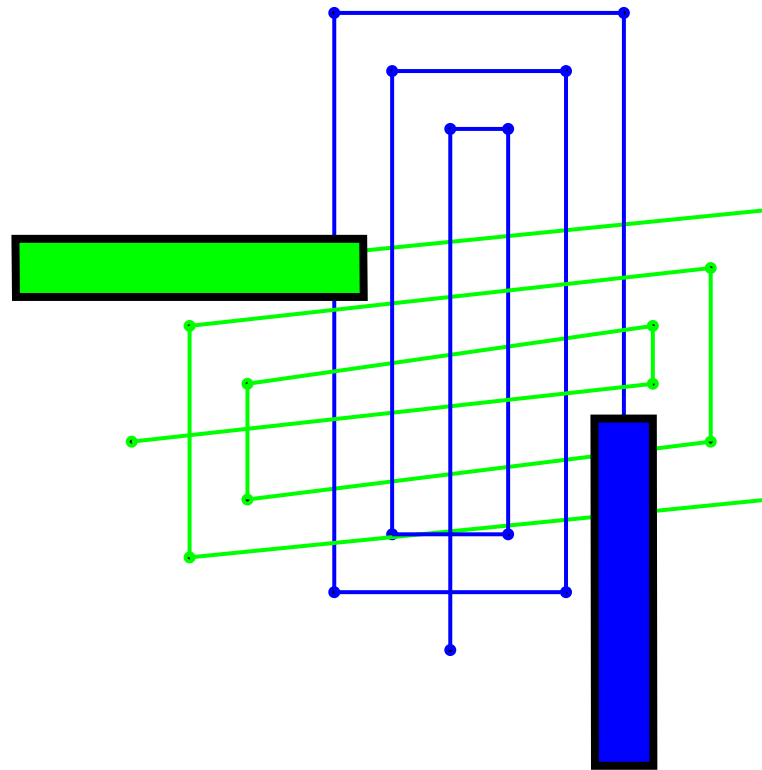
$$J(\gamma_2) = (21, 15)$$

$$J(\gamma_3) = (19, 25)$$



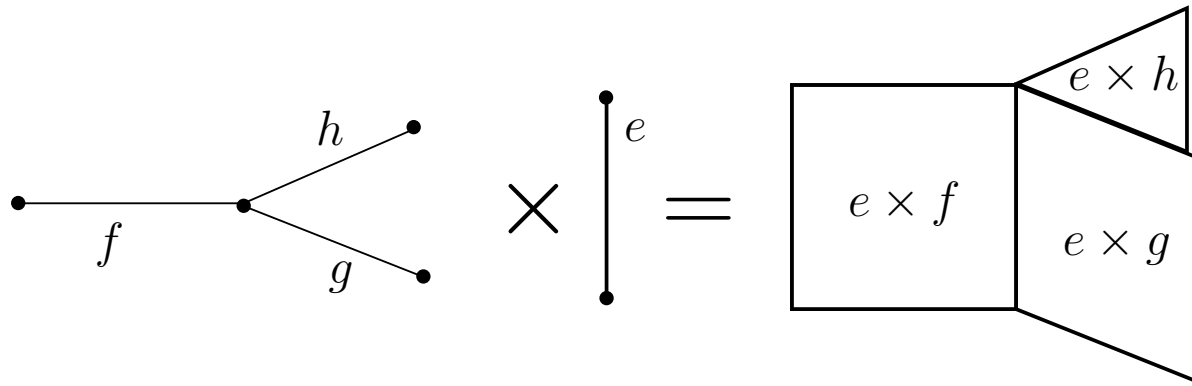
# Example

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# Acyclic Roadmaps

$\mathcal{G}_1 \times \mathcal{G}_1$  is a collection of 2-dimensional cells pasted together at their boundaries.



Same method works. Only need a technique to compute the visibility graph.

# Visibility in $\mathcal{G}_1 \times \mathcal{G}_2$

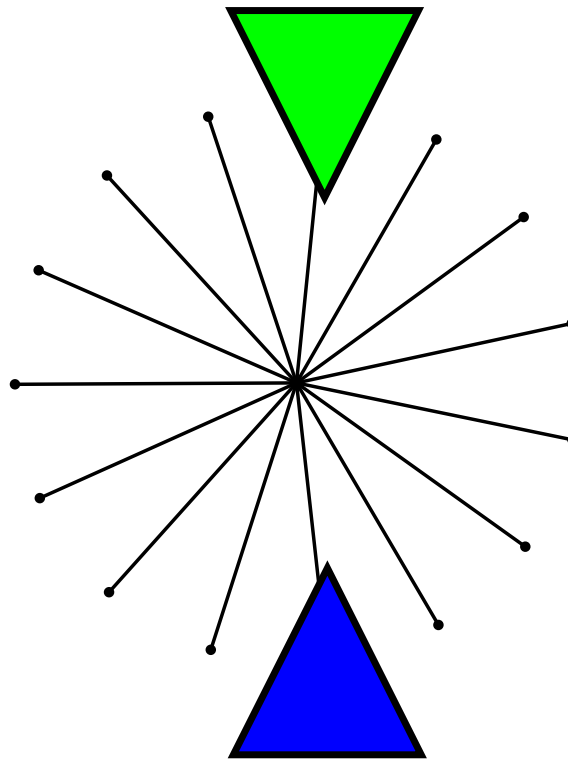
*Standard algorithm for  $\mathbb{R}^2$  (Lee, 1978):* Radial sweep about each vertex. Maintain a balanced tree of intersected segments.  $O(n^2 \log n)$  time.

*Our extension:* Radial sweep in  $\mathcal{G}_1 \times \mathcal{G}_2$ . Maintain a separate balanced tree in each for each cell.

- A ray in  $\mathcal{G}_1 \times \mathcal{G}_2$  passes through at most  $2m$  cells, where  $m$  is the total number of edges.
- At most  $2m$  binary tree operations to process each event.
- $O(mn^2 \log n)$  time.

# Example

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# Conclusion

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- Pareto optimality is an important solution concept for multiple robot coordination.
- Presented an  $O(m^2 n \log n)$  time algorithm to compute all Pareto optima for problems with  $m$  edges in the roadmaps and  $n$  obstacle vertices.
- Future Work:
  - $n$  robots. (with R. Ghrist)
  - Cyclic graphs.